



EASTERN UNIVERSITY, SRI LANKA

SPECIAL DEGREE EXAMINATION IN MATHEMATICS

(2001/2002) (Jan./Feb.2004)

MT 410-NUMERICAL LINEAR ALGEBRA

You should answer all questions. Time allowed is **THREE** hours only. Each question carries **ONE HUNDRED** marks. The numbers beside the questions indicate the approximate marks that can be gained from the corresponding parts of the questions.

1. (a) Define the terms "positive definite" and "elementary lower-triangular" as applied to an $n \times n$ Hermitian matrix A .

[10]

(b) Prove that a positive definite matrix can be expressed as $A = LU$, where L is a unit lower triangular matrix and U is an upper triangular matrix.

[25]

(c) Show that a Hermitian matrix A is positive definite if, and only if $A = GG^H$, where G is a non-singular lower triangular matrix.

[30]

Determine G such that

$$GG^H = \begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 2 & 0 & 0 \\ 1 & 0 & 4 & -2 \\ 0 & 0 & -2 & 3 \end{bmatrix}$$

[35]

2. (a) Define the terms "unitary matrix" and "elementary Hermitian matrix". [10]

(b) Show that, for any real vector x , there is a real elementary Hermitian matrix $H(w)$ such that $H(w)x = ce_1$, where $c = x^T x$ and $e_1 = (1, 0, 0, \dots, 0)^T$.

What is the optimal choice of the sign of c for the computation of w ? [30]

(c) Determine an upper triangular matrix U such that $HA = U$, where H is a product of elementary Hermitian matrices and

$$A = \begin{bmatrix} 1 & -3 & 2 \\ 2 & 4 & -1 \\ 2 & 5 & 0 \end{bmatrix}$$

making the optimal choice of sign in each stage of process. Hence solve the system $Ax = b$, where $b = (5, 0, -1)^T$ [60]

3. (a) Define the phrase strictly diagonally dominant applied to an $n \times n$ matrix A . [10]

(b) Let $A = I - L - U$ be strictly diagonally dominant, where L is strictly lower triangular and U is strictly upper triangular. For arbitrary $x^{(0)}$, a sequence $\{x^{(r)}\}$ is defined by

$$x^{(r+1)} = (I - wL)^{-1}\{wb + [(1 - w)I + wU]x^{(r)}\}, \quad r = 0, 1, 2, 3, \dots$$

Show that $x - x^{(r+1)} = M(x - x^{(r)})$, $r = 0, 1, 2, \dots$, where

$M = (I - wL)^{-1}[(1 - w)I + wU]$ and $Ax = b$. State a necessary and sufficient condition for $\{x^{(r)}\}$ to converge to x . [15]

(c) Let $0 < w \leq 1$ and let λ be any complex number with $|\lambda| \geq 1$. Show that $|\lambda + w - 1| \geq |w\lambda| \geq w$. Deduce that if λ is any eigenvalue of M , then $|\lambda| < 1$. [35]

(d) The following equations are to be solved by successive over-relaxation with a relaxation parameter 1.1.

Starting with $x^{(0)} = 0$, obtain $x^{(1)}$, $x^{(2)}$ and bound for $\|x - x^{(2)}\|_\infty$.



$$\begin{bmatrix} 11 & 1 & 0 & 0 \\ 1 & 11 & 1 & 0 \\ 0 & 1 & 11 & 2 \\ 0 & 0 & 2 & 11 \end{bmatrix} x = \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

[40]

4. (a) Define the term "upper Hessenberg matrix."

[10]

(b) i. Let A be an $n \times n$ matrix. Describe how a non-singular matrix S , a product of elementary lower triangular matrices and elementary permutation matrices, can be obtained so that $S^{-1}AS$ is an upper Hessenberg matrix.

[35]

(c) Given

$$A = \begin{bmatrix} 2 & -1 & 2 & 0 \\ 0 & -1 & 1 & 1 \\ -1 & -1 & 2 & 1 \\ -1 & 0 & -1 & 2 \end{bmatrix},$$

find an upper Hessenberg matrix $S^{-1}AS$, where S is a product of elementary permutation matrices and elementary lower triangular matrices.

[55]

5. (a) Let A be an $n \times n$ symmetric positive definite matrix. Show that the solution of the system $Ax = b$ is equivalent to the unique minimum of the function

$$F(y) := \frac{1}{2}y^T Ay - y^T b.$$

[20]

(b) For given initial iterate x_0 , the k th iterate x_k is given by

$$x_k = x_{k-1} + \alpha p_k,$$

where p_k is the search direction to be chosen such that $p_k^T r_{k-1} \neq 0$, $r_k = b - Ax_k$ is residual. Show that

$$\alpha = \frac{p_k^T r_{k-1}}{p_k^T A p_k}$$

minimizes the function $F(x_{k-1} + \alpha p_k)$ with respect to α . [20]

(c) Let A be an $n \times n$ symmetric positive definite matrix and $b \in \mathbb{R}^n$. The Conjugate Gradient (CG) iterative method for solving the system $Ax = b$ is given by, for given initial iterate $x_0 = 0$,

$$\text{Set } p_0 = r_0.$$

$$\text{While } r_k \neq 0,$$

$$\alpha_k = \frac{r_k^T r_{k-1}}{p_k^T A p_k}, \quad (\text{CG1})$$

$$x_{k+1} = x_k + \alpha_k p_k, \quad (\text{CG2})$$

$$r_{k+1} = r_k - \alpha_k A p_k, \quad (\text{CG3})$$

$$\beta_{k+1} = \frac{r_{k+1}^T r_{k+1}}{r_k^T r_k}, \quad (\text{CG4})$$

$$p_{k+1} = r_{k+1} + \beta_{k+1} p_k. \quad (\text{CG5})$$

Show that

$$\langle r_0, r_1, \dots, r_{k-1} \rangle = \langle b, Ab, \dots, A^{k-1}b \rangle.$$

[20]

Show also that

$$\begin{aligned} r_k^T r_j &= 0 \quad \text{for all } j < k \quad \text{and} \\ p_k^T A p_j &= 0 \quad \text{for all } j < k. \end{aligned}$$

[40]

(a) i. Suppose that the eigenvalue λ_1 of largest modulus and corresponding eigenvector z_1 of an $n \times n$ matrix A have been computed by the Power method. Show that there is a non-singular matrix S , a product of an elementary permutation matrix and an elementary lower triangular matrix, such that

$$A = S \left[\begin{array}{c|c} \lambda_1 & \gamma^T \\ \hline O & B \end{array} \right] S^{-1},$$



where B is an $(n - 1) \times (n - 1)$ matrix and γ is an $(n - 1)$ -column vector.

[25]

ii. Describe how the other eigenvalues and eigenvectors of A could be computed. [20]

(b) It is given that the matrix

Answer Four Questions Only

$$A = \begin{bmatrix} 2 & 1 & 0 \\ 1 & 3 & -1 \\ 0 & 1 & 0 \end{bmatrix}$$

has an eigenvalue close to 3.4 and that a corresponding eigenvector approximately $(0.7, 1, 0.3)^T$. Obtain 2×2 matrix B whose eigenvalues approximate the other eigenvalues of A . [30]

Give suitable example in each case.

(i) A large problem of size n can be divided into a sub-problems each of size n/c . The amount of time required to perform this decomposition as well as the time required to combine the solutions to the sub-problems in order to produce the final solution is $b \cdot n$. Let $T(n)$ be the running time for the problem of size n and assume $T(1)$ is constant.

Prove that

$$\begin{aligned} \text{for } n > 1 \quad T(n) &= O(n) && \text{if } a < c \\ &= O(n \log_c n) && \text{if } a = c \\ &= O(n^{\log_c a}) && \text{if } a > c \end{aligned}$$

where a , b , and c are constants.

Describe a divide-and-conquer algorithm to multiply two n -bit integers and obtain the time complexity of your method.

(ii) A single server is set up to service a sequence of customers. Assume that we know in advance how much time is required to service each customer. Use a greedy algorithm to minimize the average time spent on each customer.